DESIGN AND DEVELOPMENT OF AN ULTRA LOW POWER COMMUNICATION PROTOCOL FOR WIRELESS AUTONOMOUS MICROSYSTEMS

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TNO | Knowledge for business

Outline

Introduction

- Applications
- Communication nodes
- Protocols
- Conclusions







The next generation

Bell's law: • every decade a new generation log (people per computer) (class of computers) October 5, 2005 0 1960 1970 1980 1990 2000 2010 [Culler:2004] year

Introduction



Ultra low power protocol

- To connect autonomous micro-sensor systems
- Limited energy available
- Sensor Nodes (SN's) autonomously ACT:
 - Analyse environment
 - Communicate
 - Take action





Wireless Sensor Networks

Integrated devices

- power supply
- sensors
- embedded processor
- wireless link

Many, cheap sensors

- wireless → easy to install & operate
- intelligent → collaboration
- low-power → long lifetime

Small size

- Increased applicability
- Low production cost
- MST
- MCM (ASIC expensive)



Autonomous sensor



Transceiver

Embedded

Processor

Battery

Memory

Sensors

TNOdes

The battery crisis (Moore's law evil twin brother) Limited capacity







~2 kcal (per battery) Slow increase of capacity

- ~8% yearly increase (Wh/cm³)
- doubles every 9 years

Decrease in energy consumption is decrease in size!

~280 kcal (without cheese !)

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Sensor applications Medical

Process industry



and many more...



Urban warfare

Fire fighting

Photo: Space



Photo: RNLN

Harbor defense

(Sensor) Applications

Ad-hoc communication (no infrastructure)

- Between swarms of (micro)-robots
- Between (micro)-robots and spacecraft
- Intra-spacecraft (robust communication)



Photo: ESA



Photo: NASA



Applications @ TNO

Homecare



Munitions' health monitoring



Public safety



Agriculture



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Autonomous sensor nodes

 Pressure sensor with 868 MHz tranceiver and 2,4 Ghz power

harvesting

• Wireless microbrick sensors



TNO node (TNOde)

- Universal node
- Universal sensor interface available
 - Analog
 - Digital
- T-MAC protocol



Research topics

Issues:

- node localization
- MAC protocols
- ad-hoc routing
- network intelligence
- energy harvesting



Objectives / constraints:

- unattended operation
 - self-configuration
 - robustness
- limited resources
 - energy
 - memory



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Wireless MAC protocols

Control access to the shared medium (radio channel)

- avoid interference between transmissions
- mitigate effects of collisions (retransmit)

Approaches

- contention-based: no coordination → CSMA/CA
- schedule-based: central authority (access point) → TDMA

Traditional MAC protocols designed to:

- maximize packet throughput
- minimize latency
- provide fairness





Communication patterns

WSN applications:

- local collaboration when detecting a physical phenomenon
- periodic reporting to sink

Characteristics:

- low data rates < 1000 bps
- small messages ~ 25 bytes
- fluctuations (in time and space)





Requirements for the protocol

Handle scarce resources

- CPU: 1 10 MHz
- memory: 2 4 KB RAM
- radio: ~100 Kbps
- energy: small batteries

Unattended operation

- Plug & play
- Robustness
- Ad-hoc routing
- Long lifetime



Standards do not suffice!

- 802.11
 - ad-hoc / hopping ✓
 - power consumption (power save mode not for multi-hop networks)
 - memory footprint ^k
- Bluetooth
 - ad-hoc / hopping (limited network size)
 - power consumption [#]
 - memory footprint #
- Zigbee
 - ad-hoc / hopping (no communication between RFDs)
 - power consumption (continuous listening for peer-2-peer data transfers)

(i)

memory footprint

Design guidelines

- Switch radio off whenever possible (duty cycle) AND, minimize number of switches
- Low complexity (memory footprint)
- Minimize idle listening
- Trade off performance for energy
- Minimize overhead
- Optimize for traffic patterns

Goal:

Energy for communication
<
Energy for sensing



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Sources of overhead

- idle listening (to handle potentially incoming messages)
- collisions (wasted resources at sender and receivers)
- overhearing (communication between neighbors)
- protocol overhead (headers and signaling)
- traffic fluctuations (overprovisioning and/or collapse)
- **scalability/mobility** (additional provisions)



Dynamic duty-cycling: T-MAC



active period ends when no activation event occurs in TA

MAC protocol designed to:

- minimize energy consumption
 - prevent collisions
 - minimize protocol overhead
 - avoid overhearing
- support self-configuration



Conclusions

- Standard wireless protocols do not suffice for WSN, because of:
 - complexity (memory footprint)
 - power consumption (lifetime)
- Zigbee comes closest to being useful
 - only for particular deployment scenarios it suffices
- Solutions already exist @ TNO
 - extend Zigbee with T-MAC to truly support ad-hoc networking!



Thank you for your attention!

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END

Spare slides



Standard 1: 802.11



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Standard 2: Bluetooth



 asynchronous / synchronous / isochronous data transports



Standard 3: *Zigbee* ≈ 802.15.4





Mesh

Cluster Tree

- coordinated / uncoordinated communication
- slotted / unslotted CSMA/CA

PAN coordinatorFull Function DeviceReduced Function Device

[Zigbee Alliance:2004]

Standards overview

	Zigbee/802.15.4	Bluetooth/802.15.1	WLAN/802.11b
Application area	Monitoring	Cable replacement	E-mail, video, web
System requirements	4 - 32 kbyte	> 250 kbyte	> 1 Mbyte
Battery life (days)	100 - 1000	1 - 7	0.1 - 5
Network nodes	255 - 65k	7	30
Bandwidth (kbit/s)	250	720	>11.000
Range (m)	1-75	1-10	1 – 100
Key attributs	Cost, Power consumption	Cost, ease of usage	Fast and flexible
Source	HELICOMM		



D3: Data-centric MAC & routing

- light-weight
- scalable
- energy-conserving
- robust

No central authority

The *data* is important, not the *node* it came from





Step 1: Interest propagation



• *Sink* sends *interests* to all nodes using *flooding*





Step 2: Data advertisement



- Sink sends interests to all nodes using flooding
- Node *advertises* it *will* send data using *broadcasting*



Step 3: Data transmission



- Sink sends interests to all nodes using flooding
- Node *advertises* it *will* send data using *broadcasting*
- Actual data is *broadcasted* to *interested* neighbors





Step 4: Data advertisement revisited



- Sink sends interests to all nodes using flooding
- Node advertises it will send data using broadcasting
- Actual data is *broadcasted* to *interested* neighbors
- Subsequent advertisements are used to acknowledge and delay transmissions



Step 5: Data transmission



- Sink sends interests to all nodes using flooding
- Node advertises it will send data using broadcasting
- Actual data is *broadcasted* to *interested* neighbors
- Subsequent advertisements are used to acknowledge and delay transmissions
- The data is *forwarded* to the sink



Step 6: Final data advertisement



- Sink sends interests to all nodes using flooding
- Node advertises it will send data using broadcasting
- Actual data is *broadcasted* to *interested* neighbors
- Subsequent advertisements are used to acknowledge and delay transmissions
- The data is *forwarded* to the sink
- The sink *acknowledges* the reception of the data

Assumptions

- Transmitting data advertisements is cheaper than transmitting actual data messages
- Duty-cycled CSMA/CA MAC scheme (e.g., S-MAC, T-MAC)
- The *data* is important, not the *node* it came from

